

# The Implementation Duality

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# 1. Introduction

- Duality plays an important role in studying implementation and matching problems with quasilinear (transferable) utility.
- In the **absence of quasilinearity** much of the relevant structure is lost, but not all ....
- We exhibit a natural duality, which provides a characterization of implementability.
- We show how this characterization can be used in **matching** and **principal-agent** problems.



## 2. Model

### Basic Ingredients

- Compact metric spaces  $X$  and  $Y$ .
- $\phi : X \times Y \times \mathbb{R} \rightarrow \mathbb{R}$ , which is
  - ▶ continuous,
  - ▶ strictly decreasing in its third argument,
  - ▶ and satisfies  $\phi(x, y, \mathbb{R}) = \mathbb{R}$ .
- $\psi : Y \times X \times \mathbb{R} \rightarrow \mathbb{R}$ , which is defined as the inverse of  $\phi$  with respect to the third argument

$$u = \phi(x, y, \psi(y, x, u))$$

and inherits its properties:  $\psi$  is

- ▶ continuous,
- ▶ strictly decreasing in its third argument,
- ▶ satisfies  $\psi(y, x, \mathbb{R}) = \mathbb{R}$



## 2. Model

### Interpretation

- In the **principal-agent context**

- ▶  $\phi(x, y, v)$  is the utility of an agent of type  $x \in X$  when choosing decision  $y \in Y$  and making transfer  $v \in \mathbb{R}$  to the principal.
- ▶  $\psi(y, x, u)$  specifies the transfer that provides an agent of type  $x$  who chooses decision  $y$  with utility  $u$ .
- ▶ (we later specify a utility function for the principal, a measure over  $X$ , describing the distribution of agent types, and reservation utilities for the agent)

- In the **matching context**

- ▶  $\phi(x, y, v)$  is the maximal utility an agent of type  $x \in X$  can obtain when matched with an agent of type  $y \in Y$  who obtains utility  $v$ .
- ▶  $\psi(y, x, u)$  is the maximal utility an agent of type  $y \in Y$  can obtain when matched with an agent of type  $x \in X$  who obtains utility  $u$ .
- ▶ (we later specify measures of  $X$  and  $Y$  and reservations utilities for all agents)



## 2. Model

### Profiles and Assignments

- Let
  - ▶  $\mathbf{B}(X)$  be the set of bounded functions  $X \rightarrow \mathbb{R}$  and  $\mathbf{B}(Y)$  the set of bounded functions  $Y \rightarrow \mathbb{R}$ .
  - ▶  $Y^X$  be the set of functions  $X \rightarrow Y$  and  $X^Y$  the set of functions  $Y \rightarrow X$
- $\mathbf{u} \in \mathbf{B}(X)$  and  $\mathbf{v} \in \mathbf{B}(Y)$  are **profiles**.
- $\mathbf{y} \in Y^X$  and  $\mathbf{x} \in X^Y$  are **assignments**.
- We endow the sets  $\mathbf{B}(X)$  and  $\mathbf{B}(Y)$  with the pointwise partial order and the sup norm  $\|\cdot\|$
- (We show in the paper that the restriction to bounded profiles is without loss of generality.)



## 2. Model

### Implementation

- A profile  $\mathbf{v} \in \mathbf{B}(Y)$  implements  $(\mathbf{u}, \mathbf{y}) \in \mathbf{B}(X) \times Y^X$  if

$$\mathbf{u}(x) = \max_{y \in Y} \phi(x, y, \mathbf{v}(y))$$

$$\mathbf{y}(x) \in \arg \max_{y \in Y} \phi(x, y, \mathbf{v}(y)).$$

- Similarly, a profile  $\mathbf{u} \in \mathbf{B}(X)$  implements  $(\mathbf{v}, \mathbf{x}) \in \mathbf{B}(Y) \times X^Y$  if

$$\mathbf{v}(y) = \max_{x \in X} \psi(y, x, \mathbf{u}(x))$$

$$\mathbf{x}(y) \in \arg \max_{x \in X} \psi(y, x, \mathbf{u}(x)).$$

- We let  $I(X) \subset \mathbf{B}(X)$  and  $I(Y) \subset \mathbf{B}(Y)$  denote the sets of implementable profiles.



# 3. Duality

## Implementation Maps

- The **implementation maps**  $\Phi : \mathbf{B}(Y) \rightarrow \mathbf{B}(X)$  and  $\Psi : \mathbf{B}(X) \rightarrow \mathbf{B}(Y)$  are defined by setting

$$\Phi \mathbf{v}(x) = \sup_{y \in Y} \phi(x, y, \mathbf{v}(y)) \quad \forall x \in X$$

$$\Psi \mathbf{u}(y) = \sup_{x \in X} \psi(y, x, \mathbf{u}(x)) \quad \forall y \in Y.$$

- These implementation maps correspond to dualities in the sense of Singer (1997)



# 3. Duality

## Implementation Maps

### Proposition 1

*The images of the implementation maps coincide with the sets of implementable profiles:*

$$\mathbf{I}(X) = \Phi\mathbf{B}(Y) \text{ and } \mathbf{I}(Y) = \Psi\mathbf{B}(X).$$

- It is immediate from the definitions that implementable profiles are contained in the images of the implementation maps.
- The other direction requires more work and our assumptions on  $(X, Y, \phi)$ .





### 3. Duality

#### Galois connection

#### Proposition 2

The implementation maps  $\Phi$  and  $\Psi$  are a Galois connection. That is,

$$\mathbf{u} \geq \Phi \mathbf{v} \iff \mathbf{v} \geq \Psi \mathbf{u}$$

holds for all  $\mathbf{u} \in \mathbf{B}(X)$  and  $\mathbf{v} \in \mathbf{B}(Y)$ .

Proof:

$$\begin{aligned} \mathbf{u} \geq \Phi \mathbf{v} &\iff \mathbf{u}(x) \geq \sup_{y \in Y} \phi(x, y, \mathbf{v}(y)) \text{ for all } x \in X \\ &\iff \mathbf{u}(x) \geq \phi(x, y, \mathbf{v}(y)) \text{ for all } x \in X \text{ and } y \in Y \\ &\iff \psi(y, x, \mathbf{u}(x)) \leq \mathbf{v}(y) \text{ for all } x \in X \text{ and } y \in Y \\ &\iff \mathbf{v}(y) \geq \sup_{x \in X} \psi(y, x, \mathbf{u}(x)) \text{ for all } y \in Y \\ &\iff \mathbf{v} \geq \Psi \mathbf{u}. \end{aligned}$$



### 3. Duality

#### Galois connection

Galois connections have many nice properties. For instance:

#### Corollary 1

The implementation maps  $\Phi$  and  $\Psi$

[1.1] satisfy the **cancellation rule**, that is, for all  $\mathbf{u} \in \mathbf{B}(X)$  and  $\mathbf{v} \in \mathbf{B}(Y)$ :

$$\mathbf{v} \geq \Psi\Phi\mathbf{v} \text{ and } \mathbf{u} \geq \Phi\Psi\mathbf{u};$$

[1.2] are **order reversing**, that is, for all  $\mathbf{u}_1, \mathbf{u}_2 \in \mathbf{B}(X)$  and  $\mathbf{v}_1, \mathbf{v}_2 \in \mathbf{B}(Y)$ :

$$\mathbf{v}_1 \geq \mathbf{v}_2 \Rightarrow \Phi\mathbf{v}_2 \geq \Phi\mathbf{v}_1 \text{ and } \mathbf{u}_1 \geq \mathbf{u}_2 \Rightarrow \Psi\mathbf{u}_2 \geq \Psi\mathbf{u}_1;$$

[1.3] and satisfy the **semi-inverse rule**, that is, for all  $\mathbf{u} \in \mathbf{B}(X)$  and  $\mathbf{v} \in \mathbf{B}(Y)$ :

$$\Phi\Psi\Phi\mathbf{v} = \Phi\mathbf{v} \text{ and } \Psi\Phi\Psi\mathbf{u} = \Psi\mathbf{u}.$$



# 3. Duality

## Characterizing Implementability

Using the semi-inverse rule and Proposition 1:

### Proposition 3

[3.1]  $\mathbf{u} \in \mathbf{B}(X)$  is implementable if and only if  $\mathbf{u} = \Phi\Psi\mathbf{u}$ .

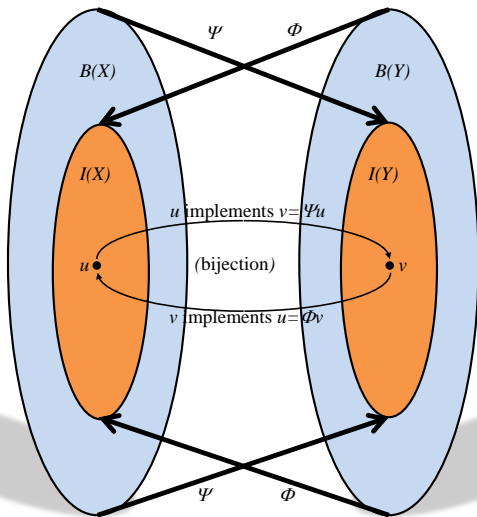
[3.2]  $\mathbf{v} \in \mathbf{B}(Y)$  is implementable if and only if  $\mathbf{v} = \Psi\Phi\mathbf{v}$ .

[3.3]  $\mathbf{u} = \Phi\mathbf{v} \Leftrightarrow \mathbf{v} = \Psi\mathbf{u}$  for all  $\mathbf{u} \in \mathbf{I}(X)$  and  $\mathbf{v} \in \mathbf{I}(Y)$ .



# 3. Duality

## Characterizing Implementability



### 3. Duality

#### Characterizing Implementability

These observations extend to assignments, showing that “everything implementable is uniquely implemented by something implementable:”

#### Corollary 2

*[2.1] If  $(\mathbf{u}, \mathbf{y}) \in \mathbf{B}(X) \times Y^X$  is implementable, then there is a unique implementable profile  $\mathbf{v}$  implementing it, namely  $\mathbf{v} = \Psi\mathbf{u}$ .*

*[2.2] If  $(\mathbf{v}, \mathbf{x}) \in \mathbf{B}(Y) \times X^Y$  is implementable, then there is a unique implementable profile  $\mathbf{u}$  implementing it, namely  $\mathbf{u} = \Phi\mathbf{v}$ .*

Hence, there is **no loss of generality in restricting attention to profiles  $\mathbf{u}$  and  $\mathbf{v}$  that implement each other.**



# 3. Duality

## Characterizing Implementability

### Corollary 3

*If  $\mathbf{u}$  and  $\mathbf{v}$  implement each other, then the argmax correspondences  $\mathbf{X}_{\mathbf{u}}$  and  $\mathbf{Y}_{\mathbf{v}}$  are inverses with*

$$\hat{x} \in \mathbf{X}_{\mathbf{u}}(\hat{y}) \Leftrightarrow \hat{y} \in \mathbf{Y}_{\mathbf{v}}(\hat{x}) \Leftrightarrow (\hat{x}, \hat{y}) \in \Gamma_{\mathbf{u}, \mathbf{v}}$$

*where*

$$\begin{aligned}\Gamma_{\mathbf{u}, \mathbf{v}} &= \{(x, y) \in X \times Y \mid \mathbf{u}(x) = \phi(x, y, \mathbf{v}(y))\} \\ &= \{(x, y) \in X \times Y \mid \mathbf{v}(y) = \psi(y, x, \mathbf{u}(x))\}\end{aligned}$$



# 3. Duality

## Some Properties of Implementable Profiles

- Implementable profiles are continuous.
- The sets of implementable profiles  $I(X)$  and  $I(Y)$  are closed subsets of  $\mathbf{B}(X)$  resp.  $\mathbf{B}(Y)$ .
- Bounded sets of implementable profiles are equicontinuous.
- Closed and bounded sets of implementable profiles are compact.



## 4. Matching

### Setting the Stage

- Matching problem is described by  $(X, Y, \phi, \mu, \nu)$ , where
  - ▶  $(X, Y, \phi)$  are as before
  - ▶  $\mu$  and  $\nu$  are measures on  $X$  and  $Y$  with full support
  - ▶  $\mu(X) = \nu(Y) > 0$
- **Remark:**
  - ▶ can append isolated types  $x_0$  and  $y_0$  to  $X$  and  $Y$  to model non-participation as a match with such a type and
  - ▶ specify measures for these types so that they absorb any excess mass on on the other side of the market.
- A **matching** for  $(X, Y, \phi, \mu, \nu)$  is a measure  $\lambda$  on  $X \times Y$  with the property that for any measurable subsets  $\tilde{X} \subset X$  and  $\tilde{Y} \subset Y$ :

$$\begin{aligned}\mu(\tilde{X}) &= \lambda(\tilde{X} \times Y) \\ \nu(\tilde{Y}) &= \lambda(X \times \tilde{Y}).\end{aligned}$$





## 4. Matching

### Pairwise Stable Outcomes

- An **outcome**  $(\lambda, \mathbf{u}, \mathbf{v})$  consists of a matching and a pair of profiles satisfying the feasibility condition  $\mathbf{u}(x) = \phi(x, y, \mathbf{v}(y))$  for all  $(x, y)$  in the support of  $\lambda$ .
- An outcome  $(\lambda, \mathbf{u}, \mathbf{v})$  is **pairwise stable** if  $\mathbf{u}(x) \geq \phi(x, y, \mathbf{v}(y))$  (or, equivalently,  $\mathbf{v}(y) \geq \psi(y, x, \mathbf{u}(x))$ ) holds for all  $(x, y) \in (X, Y)$ .

#### Lemma 1

*An outcome  $(\lambda, \mathbf{u}, \mathbf{v})$  is pairwise stable if and only if  $\mathbf{u}$  and  $\mathbf{v}$  implement each other and the support of  $\lambda$  is contained in  $\Gamma_{\mathbf{u}, \mathbf{v}}$ .*



## 4. Matching

### Existence Result

#### Proposition 4

*Every matching problem  $(X, Y, \phi, \mu, \nu)$  has a pairwise stable outcome  $(\lambda, \mathbf{u}, \mathbf{v})$ .*

**Proof** follows the same pattern as proof for existence of solution to optimal transportation problem:

- 1 Matching problems with finite numbers of agents have pairwise stable outcomes (e.g., Demange and Gale (1985))
- 2 Construct sequence of finite matching problems  $(X_n, Y_n, \phi_n, \mu_n, \nu_n)$  converging to  $(X, Y, \phi, \mu, \nu)$
- 3 Construct an associated bounded sequence of pairwise stable outcomes  $(\lambda_n, \mathbf{u}_n, \mathbf{v}_n)$
- 4 Extract converging subsequence and show that limit  $(\lambda, \mathbf{u}, \mathbf{v})$  is pairwise stable for  $(X, Y, \phi, \mu, \nu)$ .



# 4. Matching

## Stable Outcomes

- Pairwise stable outcomes can be constructed for any initial condition of the form  $\mathbf{u}(x) = \underline{u}$  for some  $x \in X$  and  $\underline{u} \in \mathbb{R}$ .
- Existence of stable outcomes then is an easy corollary to Proposition 4



# 5. Optimal Tariffs in the Principle-Agent Model

## Setting the Stage

- Agent with utility function  $\phi(x, y, v)$ .
- Principal with utility function  $\pi(x, y, v)$ .
  - ▶  $\pi : X \times Y \times \mathbb{R} \rightarrow \mathbb{R}$  is continuous, strictly increasing in  $v$  and satisfies  $\pi(x, y, \mathbb{R}) = \mathbb{R}$ .
- Agent's type distributed on  $X$  according to  $\mu$ .
- $\underline{u} \in \mathbf{B}(X)$ : reservation utility profile for the agent.



# 5. Optimal Tariffs in the Principle-Agent Model

## The Principal's Problem

Principal's problem can be formulated as:

$$\max_{(\mathbf{v}, \mathbf{u}, \mathbf{y}) \in \mathbf{I}(Y) \times \mathbf{I}(X) \times Y^X} \int_{x \in X} \pi(x, \mathbf{y}(x), \mathbf{v}(\mathbf{y}(x))) d\mu(x)$$

subject to the constraints that

- $\mathbf{v}$  implements  $(\mathbf{u}, \mathbf{y})$
- $\mathbf{u} \geq \underline{\mathbf{u}}$
- $z : X \rightarrow \mathbb{R}$ , given by  $z(x) = \pi(x, \mathbf{y}(x), \mathbf{v}(\mathbf{y}(x)))$  is measurable.



# 5. Optimal Tariffs in the Principle-Agent Model

## A Reformulation

- The incentive constraint may have multiple solutions. To address this, let

$$\hat{\pi}(x, \mathbf{v}) = \max_{y \in \mathbf{Y}_{\mathbf{v}}(x)} \pi(x, y, \mathbf{v}(y)).$$

and

$$\hat{\Pi}(\mathbf{v}) = \int_{x \in X} \hat{\pi}(x, \mathbf{v}) d\mu(x)$$

- Then the principal's problem can be written as

$$\max_{\mathbf{v} \in \mathbf{I}(Y)} \hat{\Pi}(\mathbf{v}) \text{ subject to } \mathbf{v} \leq \Psi \underline{\mathbf{u}}.$$



# 5. Optimal Tariffs in the Principle-Agent Model

## Existence Result

### Proposition 5

*A solution to the principal's problem exist.*

Proof:

- Check that  $\hat{\Pi}$  is upper semicontinuous.
- Show that there is no loss of generality in imposing a lower bound on the feasible tariffs to obtain a compact choice set.
- Apply Weierstrass.



## 6. Further Results

- Results on the lattice structure of the set of (pairwise) stable profiles carry over from the finite to the general case
  - ▶ Proof uses lattice structure induced by Galois connection in conjunction with a generalization of the [Decomposition Lemma](#)
- Sufficient conditions for the participation constraint to be binding in the principal-agent problem
  - ▶ This is a triviality with quasilinear utility
  - ▶ In general: must hold with private values, but may fail with common values





## 6. Further Results

- With  $X = [\underline{x}, \bar{x}] \subset \mathbb{R}$ ,  $Y = [\underline{y}, \bar{y}] \subset \mathbb{R}$  the **single-crossing condition**

$$\phi(x_1, y_2, v_2) \geq \phi(x_1, y_1, v_1) \Rightarrow \phi(x_2, y_2, v_2) > \phi(x_2, y_1, v_1)$$

for all  $x_1 < x_2 \in X$ ,  $y_1 < y_2 \in Y$ , and  $v_1, v_2 \in \mathbb{R}$  implies

- ▶ all increasing decision functions are implementable
- ▶ stable outcomes with deterministic matchings exist when  $\mu$  is continuous

